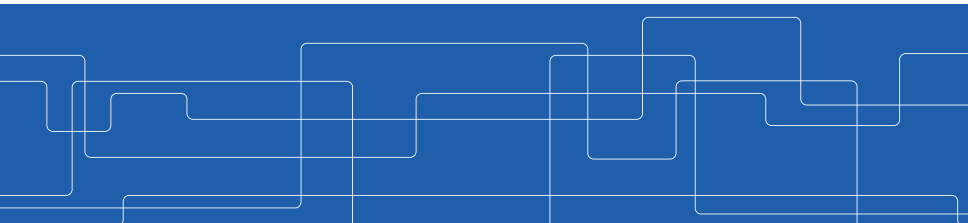


Introduction to Information Theory and its applications in Machine Learning

Amaury Gouverneur

WASP - Mathematical Foundations of AI Cluster

September 21, 2021



Outline

Shannon information content

Shannon entropy

KL-divergence

Mutual information

Application in Machine Learning

The 1948 paper

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- ▶ *A Mathematical Theory of Communication (1948)*

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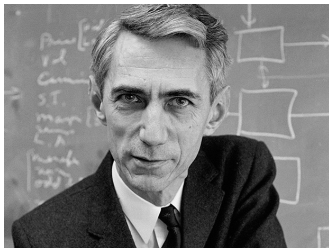


Figure: Claude Shannon

Shannon information content

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Consider a discrete random variable X .

The Shannon information content for the outcome $X = x$ is defined as:

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$$\frac{d}{dp} \log_2 \frac{1}{p} = -\frac{1}{p \ln 2} < 0 \quad \text{for } p > 0$$

Shannon information content

Verification of the properties:

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$$P(X = x, Y = y) = P(X = x)P(Y = y) \implies h(X = x, Y = y) = h(X = x) + h(Y = y)$$

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$$h(X = x) = \log_2 \frac{1}{P(X = x)} = -\log_2 P(X = x)$$

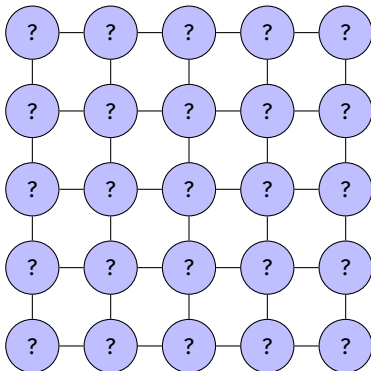
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Submarine example

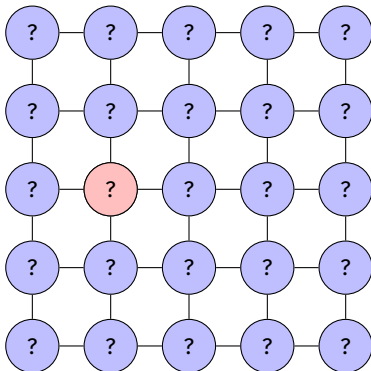
total: 0 bits



The goal is to find the submarine

Submarine example

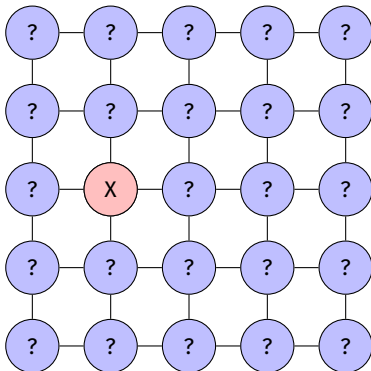
total: 0 bits



Let's chose to uncover the red "?"

Submarine example

total: 0.0588 bits

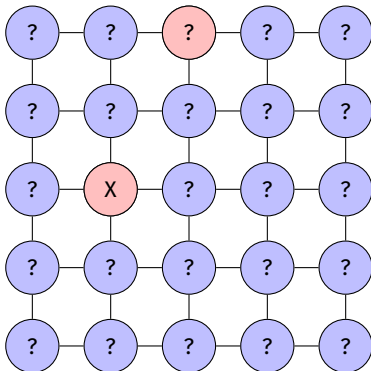


We missed. The probability to miss was $24/25$. The information we gained is

$$\begin{aligned}h(\text{miss } w/25) &= \log_2(25/24) \\ &= 0.0588\end{aligned}$$

Submarine example

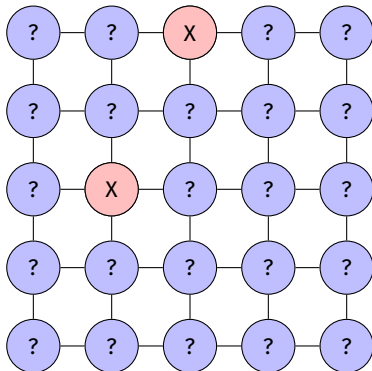
total: 0.0588 bits



Let's chose to uncover the red "?"

Submarine example

total: 0.1202 bits

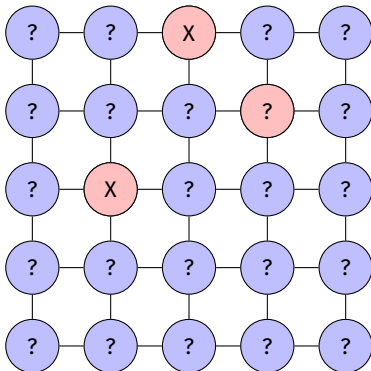


We missed. The probability to miss was $23/24$. The information we gained is

$$\begin{aligned}h(\text{miss } w/24) &= \log_2(24/23) \\ &= 0.0614\end{aligned}$$

Submarine example

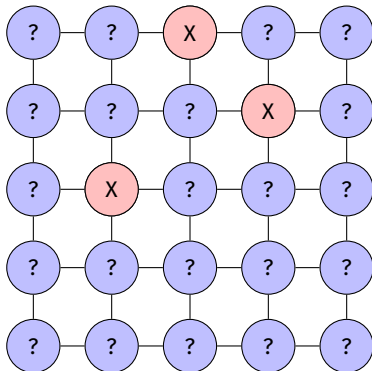
total: 0.1202 bits



Let's chose to uncover the red "?"

Submarine example

total: 0.1844 bits

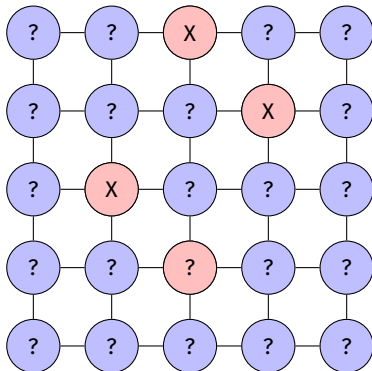


We missed. The probability to miss was $22/23$. The information we gained is

$$\begin{aligned}h(\text{miss } w/23) &= \log_2(23/22) \\ &= 0.0641\end{aligned}$$

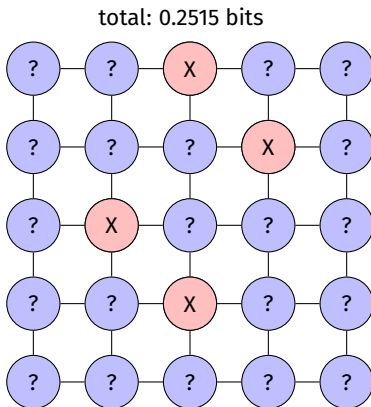
Submarine example

total: 0.1844 bits



Let's chose to uncover the red "?"

Submarine example

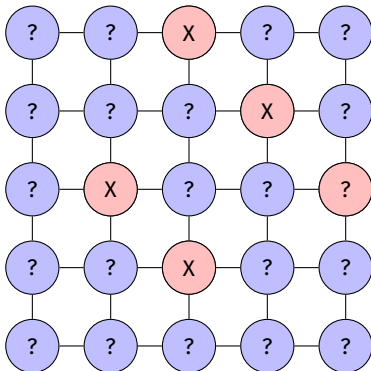


We missed. The probability to miss was $21/22$. The information we gained is

$$\begin{aligned}h(\text{miss } w/22) &= \log_2(22/21) \\ &= 0.0671\end{aligned}$$

Submarine example

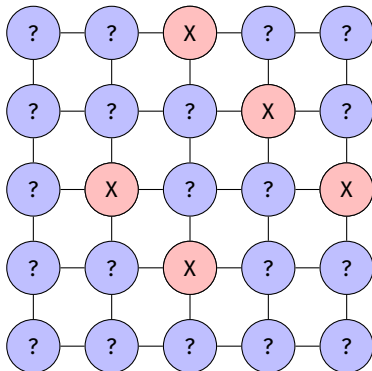
total: 0.2515 bits



Let's chose to uncover the red "?"

Submarine example

total: 0.3219 bits

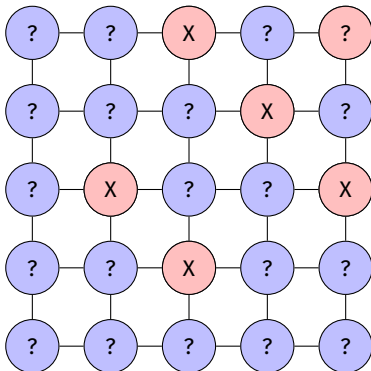


We missed. The probability to miss was $20/21$. The information we gained is

$$\begin{aligned}h(\text{miss } w/21) &= \log_2(21/20) \\ &= 0.0703\end{aligned}$$

Submarine example

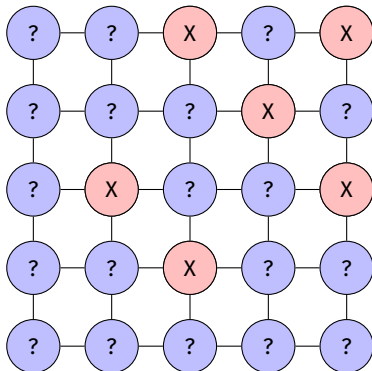
total: 0.3219 bits



Let's chose to uncover the red "?"

Submarine example

total: 0.3959 bits

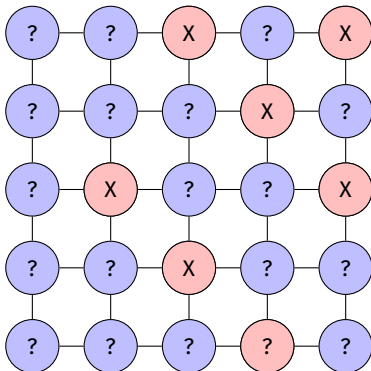


We missed. The probability to miss was $19/20$. The information we gained is

$$\begin{aligned}h(\text{miss } w/20) &= \log_2(20/19) \\ &= 0.0740\end{aligned}$$

Submarine example

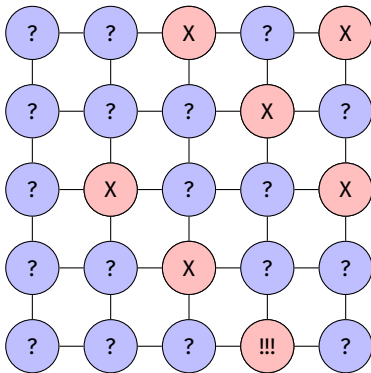
total: 0.3959 bits



Let's chose to uncover the red "?"

Submarine example

total: 4.6439 bits

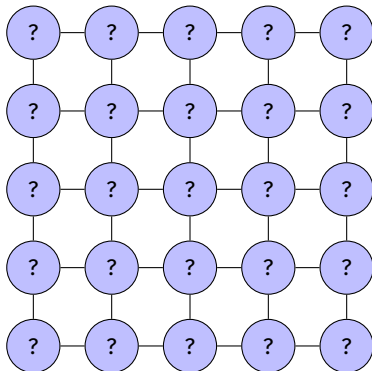


We found it! The probability to hit was $1/19$. The information we gained is

$$\begin{aligned}h(\text{hit } w/22) &= \log_2(19/1) \\ &= 4.248\end{aligned}$$

Another try

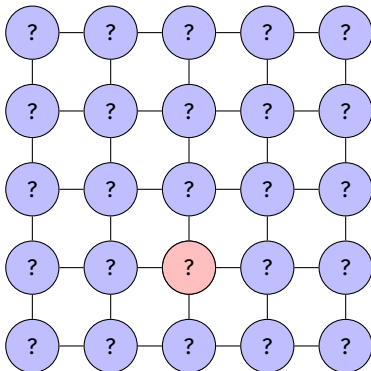
total: 0 bits



Let's try again

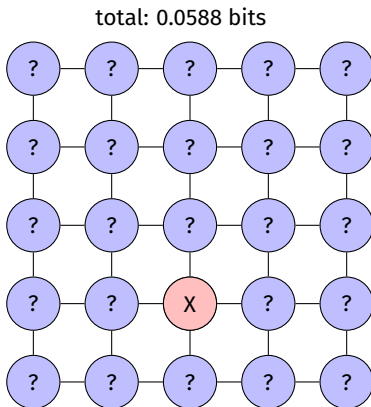
Another try

total: 0 bits



Let's chose to uncover the red "?"

Another try

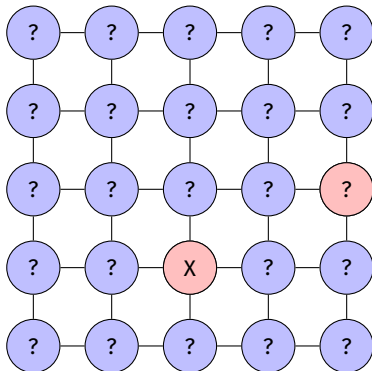


We missed. The probability to miss was $24/25$. The information we gained is

$$\begin{aligned}h(\text{miss } w/25) &= \log_2(25/24) \\ &= 0.0588\end{aligned}$$

Another try

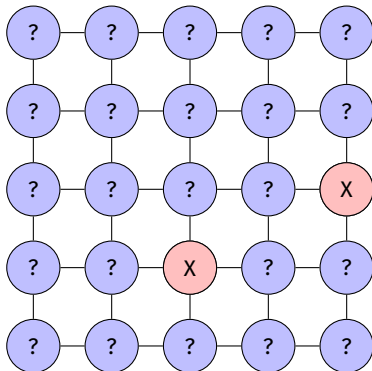
total: 0.0588 bits



Let's chose to uncover the red "?"

Another try

total: 0.1202 bits

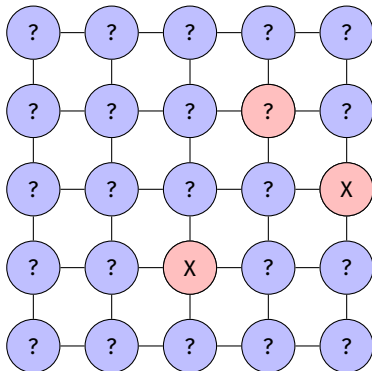


We missed. The probability to miss was $23/24$. The information we gained is

$$\begin{aligned}h(\text{miss } w/24) &= \log_2(24/23) \\ &= 0.0614\end{aligned}$$

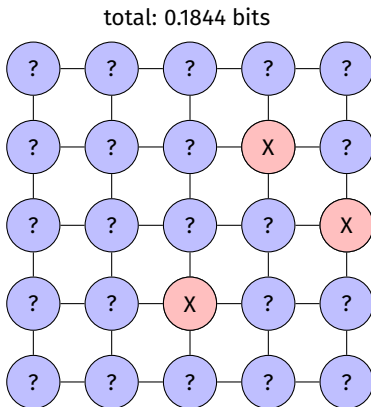
Another try

total: 0.1202 bits



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Another try

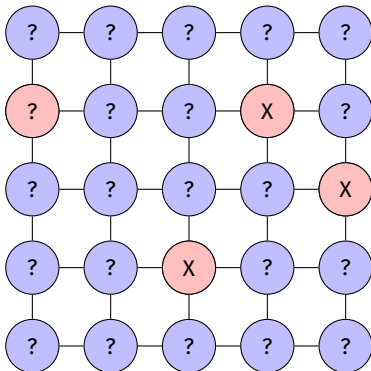


We missed. The probability to miss was $22/23$. The information we gained is

$$\begin{aligned}h(\text{miss } w/23) &= \log_2(23/22) \\ &= 0.0641\end{aligned}$$

Another try

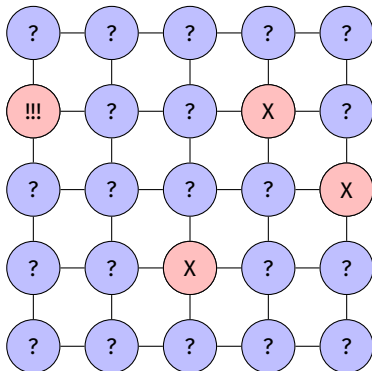
total: 0.1844 bits



Let's chose to uncover the red "?"

Another try

total: 4.6439 bits

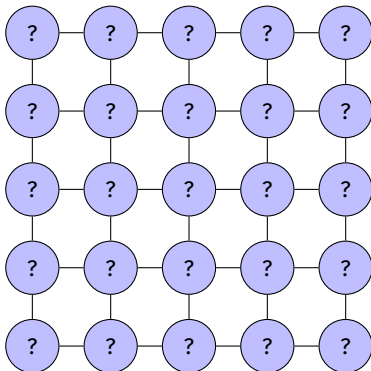


We found it! The probability to hit was $1/22$. The information we gained is

$$\begin{aligned}h(\text{hit } w/22) &= \log_2(22/1) \\ &= 4.4594\end{aligned}$$

One more try

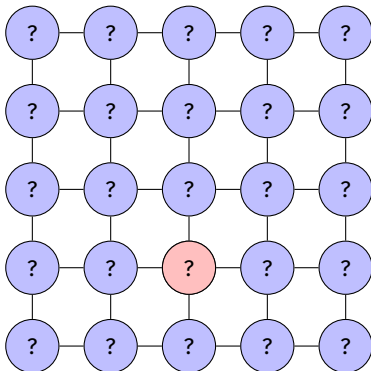
total: 0 bits



Let's try again

One more try

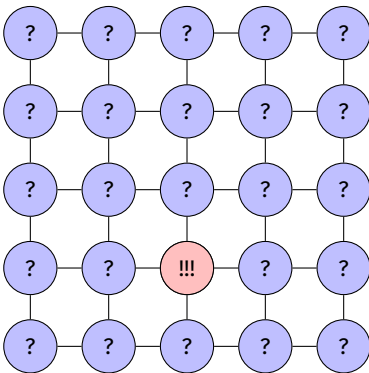
total: 0 bits



Let's chose to uncover the "?"

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We found it! The probability to hit was $1/25$. The information we gained is

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Shannon entropy

Average information content:

$$H(X) = \sum_{x \in \mathcal{X}} P(X = x) h(x = X)$$

Shannon entropy

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$$\begin{aligned} H(X) &= \sum_{x \in \mathcal{X}} P(X = x) h(x = X) \\ &= \sum_{x \in \mathcal{X}} P(X = x) \log_2 \left(\frac{1}{P(X = x)} \right) \end{aligned}$$

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Shannon entropy

Example for a weighted coin

Let

$$X = \begin{cases} 1 & \text{with probability } p \\ 0 & \text{with probability } 1 - p \end{cases}$$

Shannon entropy

Example for a weighted coin

Let

$$X = \begin{cases} 1 & \text{with probability } p \\ 0 & \text{with probability } 1 - p \end{cases}$$

Then

$$H(X) = -p \log(p) - (1 - p) \log(1 - p)$$

Shannon entropy

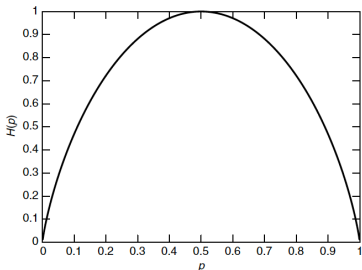
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Joint entropy

Multivariate generalization of the Shannon entropy.

$$H(X, Y) = \sum_{x \in \mathcal{X}} \sum_{y \in \mathcal{Y}} P(X = x, Y = y) \log_2 \left(\frac{1}{P(X = x)P(Y = y)} \right)$$

Conditional entropy

Entropy of the conditional distribution

$$H(X|Y = y) = \sum_{x \in \mathcal{X}} P(X = x|Y = y) \log_2\left(\frac{1}{P(X = x|Y = y)}\right)$$

Conditional entropy

Entropy of the conditional distribution

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Conditional entropy

$$H(X|Y) = \sum_{y \in \mathcal{Y}} P(Y = y)H(X|Y = y)$$

Conditional entropy

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Conditional entropy

Properties of the conditional entropy:

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1. $H(X,Y) = H(X|Y) + H(Y)$

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Conditional entropy

Properties of the conditional entropy:

1. $H(X,Y) = H(X|Y) + H(Y)$
2. $H(X|Y) = 0$ if X is deterministic knowing Y
3. $H(X|Y) = H(X)$ if X and Y are independent

Kullback-Leibler divergence

A useful "measure" of difference between two distributions. Let P and Q be two distributions,

$$D_{KL}(P||Q) = \sum_{x \in \mathcal{X}} P(x) \log_2 \frac{P(x)}{Q(x)}$$

Kullback-Leibler divergence

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Properties of the KL-divergence:

Kullback-Leibler divergence

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Properties of the KL-divergence:

1. $D_{KL}(P||Q) \geq 0$

Kullback-Leibler divergence

A useful "measure" of difference between two distributions. Let P and Q be two distributions,

$$D_{KL}(P||Q) = \sum_{x \in \mathcal{X}} P(x) \log_2 \frac{P(x)}{Q(x)}$$

Properties of the KL-divergence:

1. $D_{KL}(P||Q) \geq 0$
2. $D_{KL}(P||Q) = 0$ only if $P = Q$

Kullback-Leibler divergence

A useful "measure" of difference between two distributions. Let P and Q be two distributions,

$$D_{KL}(P||Q) = \sum_{x \in \mathcal{X}} P(x) \log_2 \frac{P(x)}{Q(x)}$$

Properties of the KL-divergence:

1. $D_{KL}(P||Q) \geq 0$
2. $D_{KL}(P||Q) = 0$ only if $P = Q$
3. It's not a metric: $D_{KL}(P||Q) \neq D_{KL}(Q||P)$ in general

KL-divergence

Illustration



KL video

Mutual information

$$I(X; Y) = D_{\text{KL}}(P(x, y) || P(x)P(y))$$

Mutual information

$$\begin{aligned} I(X; Y) &= D_{\text{KL}}(P(x, y) || P(x)P(y)) \\ &= \sum_{x \in \mathcal{X}} \sum_{y \in \mathcal{Y}} P(X = x, Y = y) \log_2 \left(\frac{P(X = x, Y = y)}{P(X = x)P(Y = y)} \right) \end{aligned}$$

Mutual information

$$\begin{aligned} I(X; Y) &= D_{\text{KL}}(P(x, y) || P(x)P(y)) \\ &= \sum_{x \in \mathcal{X}} \sum_{y \in \mathcal{Y}} P(X = x, Y = y) \log_2 \left(\frac{P(X = x, Y = y)}{P(X = x)P(Y = y)} \right) \\ &= H(X) - H(X|Y) \end{aligned}$$

Mutual information

$$\begin{aligned} I(X; Y) &= D_{\text{KL}}(P(x, y) || P(x)P(y)) \\ &= \sum_{x \in \mathcal{X}} \sum_{y \in \mathcal{Y}} P(X = x, Y = y) \log_2 \left(\frac{P(X = x, Y = y)}{P(X = x)P(Y = y)} \right) \\ &= H(X) - H(X|Y) \\ &= H(Y) - H(Y|X) \end{aligned}$$

Mutual information

$$\begin{aligned} I(X; Y) &= D_{\text{KL}}(P(x, y) || P(x)P(y)) \\ &= \sum_{x \in \mathcal{X}} \sum_{y \in \mathcal{Y}} P(X = x, Y = y) \log_2 \left(\frac{P(X = x, Y = y)}{P(X = x)P(Y = y)} \right) \\ &= H(X) - H(X|Y) \\ &= H(Y) - H(Y|X) \\ &= H(X, Y) - H(X|Y) - H(Y|X) \end{aligned}$$

Mutual information

$$\begin{aligned}I(X; Y) &= D_{\text{KL}}(P(x, y) || P(x)P(y)) \\&= \sum_{x \in \mathcal{X}} \sum_{y \in \mathcal{Y}} P(X = x, Y = y) \log_2 \left(\frac{P(X = x, Y = y)}{P(X = x)P(Y = y)} \right) \\&= H(X) - H(X|Y) \\&= H(Y) - H(Y|X) \\&= H(X, Y) - H(X|Y) - H(Y|X)\end{aligned}$$

Data-processing inequality.

Mutual information

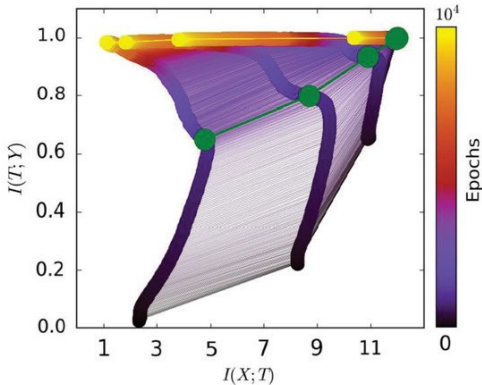
$$\begin{aligned} I(X; Y) &= D_{\text{KL}}(P(x, y) || P(x)P(y)) \\ &= \sum_{x \in \mathcal{X}} \sum_{y \in \mathcal{Y}} P(X = x, Y = y) \log_2 \left(\frac{P(X = x, Y = y)}{P(X = x)P(Y = y)} \right) \\ &= H(X) - H(X|Y) \\ &= H(Y) - H(Y|X) \\ &= H(X, Y) - H(X|Y) - H(Y|X) \end{aligned}$$

Data-processing inequality. If $X \rightarrow Y \rightarrow Z$ forms a Markov chain,

$$I(X; Y) \geq I(X; Z)$$

Application in Machine Learning

The information plane and the information bottleneck. *Shwartz-Ziv, Ravid, and Naftali Tishby. "Opening the black box of deep neural networks via information." (2017)*



DNN as an encoder-decoder

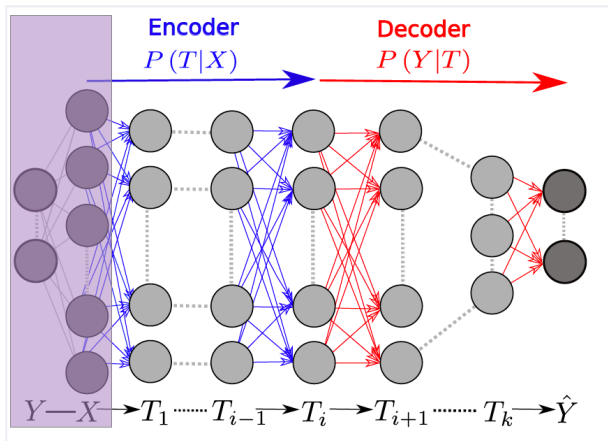


Figure: The DNN layers form a Markov chain of successive internal representations of the input layer X .

The Information Plane

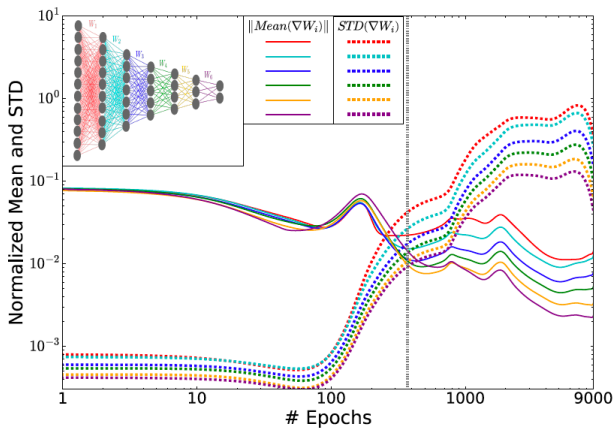
First optimization phase

Learning part 1

Second optimization phase

Learning part 2

The drift and diffusion phases of SGD optimization



The end

Thank you for listening. Any **questions**?